## From Templates to Folds

## CS 5010 Program Design Paradigms "Bootcamp"

## Lesson 6.3

## Introduction

- Last week, we saw how the built-in mapping functions on lists, like map, filter, and foldr, made writing functions on lists easier.
- In this lesson we'll see how we can do something similar for any recursive data definition.


## Learning Objectives

- At the end of this lesson you should be able to:
- Write a fold function for any recursive data definition
- Use the fold function to define useful functions on that data


## Binary Trees

(define-struct leaf (datum))
(define-struct node (lson rson))
; ; A Tree is either
; ; -- (make-node Tree Tree)

Here is the definition of a binary tree again.

## Template

tree-fn : Tree -> ? ? ?
(define (tree-fn t)
(cond
[(leaf? t) (... (leaf-datum t))]
[else (...
(tree-fn (node-lson t))
(tree-fn (node-rson t)))])
Self-reference in the data definition leads to self-reference in the template; Self-reference in the template leads to self-reference in the code.

## The template has two blanks

tree-fn : Tree -> ?? ?
(define (tree-fn t)
Two blanks: one blue and one orange
(cond
[(leaf? t) (. (leaf-datum t))]
[else
(tree-fn (node-lson t))
(tree-fn (node-rson t) ))])

## From templates to folds

- Observe that the template has two blanks: the blue one and the orange one.
- Any two functions that follow the template will be the same except for what goes in the blanks.
- So we can generalize them by adding arguments for each blank.


## Template $\rightarrow$ tree-fold

tree-fold : ... Tree -> ???
(define (tree-fold
(cond
[(leaf? t) (base (leaf-datum t))]
[else (combiner
(tree-fold
(node-lson $t$ ))
(tree-fold
(node-rson t)))])
Corresponding to each blank, we add an extra argument: combiner (in blue) for the blue blank and base (in orange) for the orange blank, and we pass these arguments to each of the recursive calls, just like we did for lists. The strategy for tree-fold is "Use template for Tree on t"

## What's the contract for tree-fold?

tree-fold

(define (tree-fold combiner base t) (cond
[(leaf? t) (base (leaf-datum t))]
[else (combiner $\leftarrow(x \times->x)$
$x \longrightarrow x$ (tree-fold combiner base (node-lson t))
(tree-fold combiner base

Since tree-fold returns an $\mathbf{X}$, the arguments to combiner are both $\mathbf{X}$ 's, and combiner itself must return an $\mathbf{X}$. (node-rson t)))]))

So combiner must be an
( X X -> X)

## Be sure to reconstruct the original functions!

(define (tree-sum t) (tree-fold + (lambda (n) n) t))
(define (tree-min $t$ )
(tree-fold min (lambda (n) n) t))

> Here are our original functions, sum, tree-min, and tree-max, rewritten using tree-fold.

The strategy for each of these is "Call a more general function."

## Another example of trees: Ancestor Trees

(define-struct person (name father mother)) (define-struct adam ()) (define-struct eve ())

```
;; A Person is either
;; -- (make-adam)
;; -- (make-eve)
;; -- (make-person String Person Person)
;; person-fn : Person
(define (person-fn p)
    (cond
    [(adam? p) ...]
    [(eve? p) ...]
        [else (...
    (person-name p)
    (person-fn (person-father p))
    (person-fn (person-mother p)))]))
```

Self-reference in the data definition leads to selfreference in the template; Self-reference in the template leads to selfreference in the code.

## Template for Person

; ; person-fn : Person -> ???
(define (person-fn p)
(cond
[(adam? p) ...]
[(eve? p) ...]
[else (
(person-name p)
(person-fn (person-father p))
(person-fn (person-mother p)))]))

## From template to fold:

; ; person-fold : ... Person -> ?? ?
(define (person-fold adam-val eve-val combiner $p$ ) (cond
[(adam? p) adam-val]
[(eve? p) eve-val]
[else (combiner
Corresponding to our three blanks we add three arguments: the value for adam (in blue), the value for eve (in purple) and the combiner (in orange).
(person-name p)
(person-fold adam-val eve-val combiner (person-father p))
(person-fold adam-val eve-val combiner (person-mother p)))]))

## What's the contract for person-fold?

## ; ; person-fold

We can work out the contract for person-fold the same way that we did for treefold. Here again we've marked some of the sub-expressions with the kind of value they return.
; ; : X X (String $\mathrm{X} \times \mathrm{X}$-> X ) Person -X
(define (person-fold adam-val eve-val combiner p)
(cond
[(adam? p) adam-val]
[(eve? p) eve-val]
x [else ${ }^{\text {(combiner }}$
(String X X $\rightarrow$ X)


Observe, as before, that the arguments to combiner match combiner's contract, and that all three branches of the cond return an $\mathbf{X}$, so the whole function is guaranteed to return an $\mathbf{X}$.

## Summary

- You should be able to:
- Write a fold function for any recursive data definition
- Use the fold function to define useful functions on that data


## Next Steps

- Study the file 06-3-tree-folds.rkt in the Examples folder.
- If you have questions about this lesson, ask them on the Discussion Board
- Do Guided Practice 6.3
- Go on to the next lesson

